

# Numeric Truncation Security Predicate

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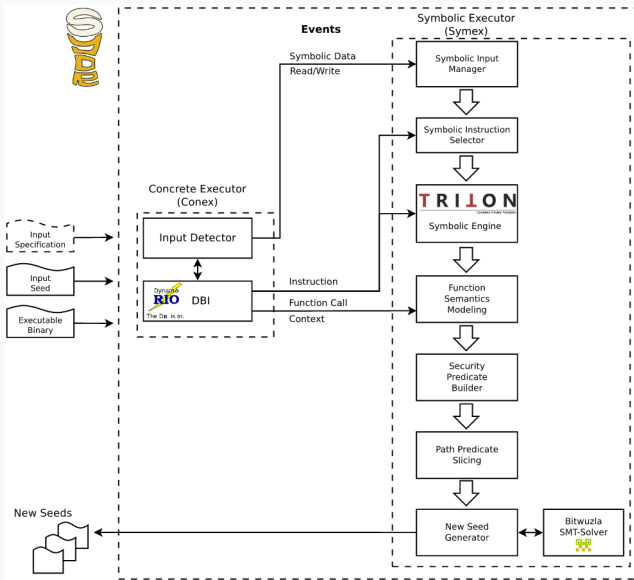
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ISP RAS

- Bugs and vulnerabilities may emerge during the development lifecycle.
- Lots of companies integrate security development lifecycle (SDL) into their workflow processes to detect errors.
- One of the most popular SDL technology is dynamic symbolic execution (DSE).

# Dynamic Symbolic Execution

- Dynamic Symbolic Execution is used for analyzing the program at runtime by constructing formulas over symbolic variables.
- These formulas can be analyzed to obtain program execution behavior.
- It is possible to find bugs in binary code with DSE building formulas corresponding to error conditions, which we call security predicates.



# Sydr and Security Predicates



Sydr uses [Triton](#) as DSE engine, [DynamoRIO](#) to instrument binary code instructions and [Bitwuzla](#) as SMT-solver.

Dynamic symbolic execution:

- Each input byte is modeled by a free symbolic variable.
- Instruction interpretation produce SMT formulas.
- Path predicate contains taken branch constraints.
- Sydr inverts branches to explore new paths and solves security predicates (integer overflow, null pointer dereference, etc.).
- Sydr sends security predicate to Bitwuzla which generates the input file to reproduce the error in case of successful solution of the constructed formula.

# Numeric Truncation

- Numeric truncation error occurs when a value with the bigger type size is converted to the smaller type.
- This error is typical for such programming languages as C/C++, Java, etc.
- Numeric truncation can lead to incorrect program execution or even to some vulnerabilities.
- For example, truncated value of allocated memory size may later cause memory corruptions; or truncated value in loop condition may lead to infinite loop.

# Simple example of Numeric Truncation

- 64-bit program
- Input: +00000065536
- Numeric Truncation in line 8
- Output: +00000000000

```
1      #include <stdio.h>
2      #include <stdint.h>
3
4      int
5      main() {
6          uint32_t a = 0;
7          scanf("%u", &a);
8          uint16_t b = a;
9          printf("%u\n", b);
10         return 0;
11     }
```

# Basic Algorithm

- We analyze every symbolic instruction during the symbolic execution.
- When we meet instructions like `mov`, `movsx`, `movzx`, `cbw`, `cwde`, `cdqe`, we check them for the numeric truncation error by building and solving security predicate.
- For `mov*` instructions: if the initial size of the symbolic value, located in register or memory, is bigger than the source operand size, then we check the security predicate.
- For convert instructions: if the size of the symbolic value located in `rax` register is bigger than the size of the value being extended after `cbw`, `cwde`, `cdqe` instructions execution, then numeric truncation error may occur.



# Security Predicate formulas

Unsigned truncation example:

0	0	0	0	0	0	1	1	1	1	1	1	0	0	0	0		1	1	1	1	1	0	0	0	1	1	1	1	1	1	1	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	--	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

Signed truncation examples:

0	0	0	0	0	0	1	1	1	0	0	0	0	0	0	0		1	1	1	1	1	0	0	0	1	1	1	1	1	1	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	--	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

1	1	1	1	1	1	0	0	0	1	1	1	1	1	1	1		1	1	1	1	1	0	0	0	1	1	1	1	1	1	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	--	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

False signed truncation examples:

0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		1	1	1	1	1	0	0	0	1	1	1	1	1	1	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	--	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1		1	1	1	1	1	1	1	1	1	1	1	1	0	1	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	--	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

Formula in security predicate:

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Unsigned formula:	Signed formula:
$\text{not}(\phi_{\text{trunc}} == \text{bv}(0, \text{sz}))$	$\text{not}((\phi_{\text{trunc}} == \text{bv}(0, \text{sz})) \vee (\phi_{\text{trunc}} == \text{bv}(1, \text{sz})))$

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$\phi_{\text{trunc}}$  — formula of bits being truncated

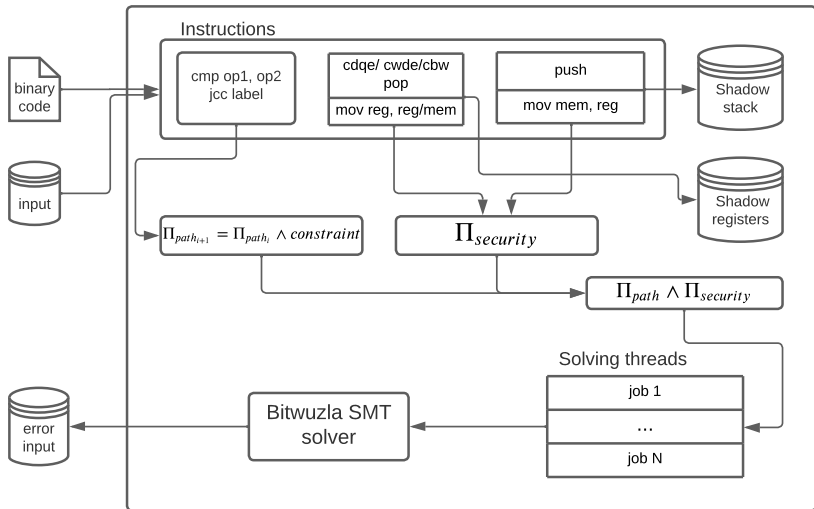
$\text{bv}(0, \text{sz})$  — bitvector of  $\text{sz}$  zeros,  $\text{bv}(1, \text{sz})$  — bitvector of  $\text{sz}$  ones

- We get the signedness when the value was read with functions like `fscanf`, `scanf`, etc., that use `strto*` functions inside.
- Otherwise, we get it with the backward slicing algorithm. We find branching instructions like `j1`, `ja`, `jb`, etc. to guess signedness.
- If we couldn't get the signedness, we build security predicate with conjunction of signed and unsigned formulas.

# False Positive Example

```
1  #include <stdio.h>
2  #include <stdint.h>
3
4  void
5  foo(int16_t a, int8_t b)
6  {
7      printf("%d%d", a, b);
8  }
9
10 int
11 main()
12 {
13     int32_t i32;
14     scanf("%d", &i32);
15     foo(i32, i32);
16     return 0;
17 }
18
19 00001255 <main>:
20  call    10c0 <__isoc99_scanf@plt>
21  add     esp,0x10
22  mov     eax,DWORD PTR [ebp-0x10]
23  movsx  edx,al
24  mov     eax,DWORD PTR [ebp-0x10]
25  cwde
26  push   edx
27  push   eax
28  call   120d <foo>
29  ...
30 0000120d <foo>:
31  ...
32  mov     edx,DWORD PTR [ebp+0x8]
33  mov     ecx,DWORD PTR [ebp+0xc]
34  mov     WORD PTR [ebp-0xc],dx
35  mov     edx,ecx
36  mov     BYTE PTR [ebp-0x10],dl
```

# Full Scheme (implementation)



- To save the size of symbolic memory operands we have Symbolic Shadow Stack. To fill the shadow stack we analyze `mov`, `movsx`, `movzx`, `push` instructions.
- For `mov`, `movsx`, `movzx` instructions:
  1. Check if the destination operand is memory and source operand is a symbolic register.
  2. Check if the source register is in shadow register map.
  3. If so, fill the shadow stack with minimum of source operand size and the size from shadow registers. Otherwise, fill it with the size of source operand.
- For `push` we search the source operand register in shadow register map, and then we fill the current `sp` value address in shadow stack the same way as for `mov`, `movsx`, `movzx` instructions.

# Shadow Registers

- To keep up-to-date the size of symbolic register operands we have Symbolic Shadow Registers. To fill the shadow registers we analyze `mov`, `movsx`, `movzx`, `cbw`, `cwde`, `cdqe` and `pop` instructions.
- For `mov`, `movsx`, `movzx` instructions:
  1. Check if destination operand is register and source operand is symbolic.
  2. Check if the source register is in shadow register map (shadow stack).
  3. If so, fill the shadow register with minimum of source operand size and the size from shadow registers (shadow stack). Otherwise, fill it with the size of source operand.
- In case of `pop` instruction we try to get symbolic size from shadow stack for current `sp` value and then save it in shadow registers map for the register operand.

# Shadow Registers


- For `cbw`, `cwde`, `cdqe` instructions we update the shadow registers map with the size being extended.
- In case of other instructions we get all the registers written by the instruction and update the shadow registers map with the size of written value. This is necessary to keep up-to-date actual symbolic sizes when the arithmetic operations are performed.

- In previous works we have adapted Juliet test suite to make it suitable for dynamic analysis. It builds all the tests for specified CWE to binaries in 32-bit and 64-bit modes. Then it runs the tool under test on all the binaries with the sample input data.
- We tested our numeric truncation security predicate on Juliet Dynamic CWE-197 and our approach showed 100% accuracy.



Sydr has found 12 new numeric truncation errors in open-source projects. All of them were reported and fixed.

Project	Detected errors number
nDPI	7
libpcap	2
FreelImage	1
LibTIFF	1
unbound	1

 — approved by maintainers

 — not approved by maintainers

```
1  struct diameter_header_t {
2      u_int8_t  com_code[3];
3      ...
4  };
5
6  typedef enum {
7      AC = 271,
8      AS = 274,
9      CC = 272,
10     CE = 257,
11     DW = 280,
12     DP = 282,
13     RA = 258,
14     ST = 275
15 } com_type_t;
16 ...
```

```
1  u_int16_t com_code =
2  diameter->com_code[2]
3  + (diameter->com_code[1] << 8)
4  + (diameter->com_code[0] << 8);
5
6  if (com_code == AC ||
7      com_code == AS ||
8      com_code == CC ||
9      com_code == CE ||
10     com_code == DW ||
11     com_code == DP ||
12     com_code == RA ||
13     com_code == ST)
14  return 0;
```

# LibTIFF

```
1 static void TIFFReadDirectoryCheckOrder(  
2     TIFF *tif, TIFFDirEntry *dir,  
3     uint16_t dircount)  
4 {  
5     static const char module[]  
6     = "TIFFReadDirectoryCheckOrder";  
7     uint16_t m;  
8     uint16_t n;  
9     TIFFDirEntry *o;  
10    m = 0;  
11    for (n = 0, o = dir;  
12         n < dircount;  
13         n++, o++)  
14    {  
15        ...  
16    }  
17 }
```

```
1     if (o->tdir_tag < m)  
2     {  
3         TIFFWarningExtR(  
4             tif, module,  
5             "Invalid TIFF directory;"  
6             "tags are not sorted in "  
7             "ascending order");  
8         break;  
9     }  
10    m = o->tdir_tag + 1;
```

```
1      uint16_t size;
2      ...
3      size = tlv->tlv_length;
4      if (size % 4 != 0)
5          size += 4 - size % 4;
6
7      if (size < sizeof(nflog_tlv_t)) {
8          /* Yes. Give up now. */
9          return;
10     }
11
12     if (caplen < size || length < size) {
13         /* No. */
14         return;
15     }
16     ...
```

**Questions?**

# Security Predicate formulas

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**Unsigned formula:**

$\text{not}(\phi_{trunc} == \text{bv}(0, sz))$

**Signed formula:**

$\text{not}((\phi_{trunc} == \text{bv}(0, sz)) \vee (\phi_{trunc} == \text{bv}(1, sz)))$

---

$\phi_{trunc} = \text{extract}(high - 1, low, \phi_{var})$  — significant bits being truncated from the original value

$high$  — original symbolic size,  $low$  — size of resulting value,  $\phi_{var}$  — symbolic variable formula

$\text{bv}(0, sz)$  — bitvector of  $sz$  zeros,  $\text{bv}(1, sz)$  — bitvector of  $sz$  ones

$sz = high + low - 1$  — size of  $\phi_{trunc}$  formula

# nDPI

```
1 void ndpi_data_add_value(  
2     struct ndpi_analyze_struct *s,  
3     const u_int64_t value) {  
4     if(!s)  
5         return;  
6     if(s->sum_total == 0)  
7         s->min_val = s->max_val = value;  
8     else {  
9         if(value < s->min_val)  
10            s->min_val = value;  
11        if(value > s->max_val)  
12            s->max_val = value;  
13    }  
14    s->sum_total += value,  
15    s->num_data_entries++;  
16    if(s->num_values_array_len) {  
17        s->values[s->next_value_insert_index]  
18            = value;  
19    }  
20    ...  
21 }
```

# unbound

```
1  int sldns_str2wire_type_buf(const char * str,
2  uint8_t* rd, size_t* len)
3  {
4  uint16_t t = sldns_get_rr_type_by_name(str);
5  ...
6  }
7
8  sldns_rr_type
9  sldns_get_rr_type_by_name(const char *name)
10 {
11 ...
12 if (strlen(name) > 4 && strncasecmp(name,
13 "TYPE", 4) == 0) {
14     return atoi(name + 4);
15 }
16 ...
17 }
```